

# Modeling Public Key Infrastructures in the Real World

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# Making Trust Judgements

- PKIs give users information to make *trust judgements*
- Based on initial assumptions and a pile of certificates
- If PKI works, we can deduce what we should and can't deduce what we shouldn't
- Complex and important decision: use formal methods
- PKI designers can verify their designs

# Maurer's Deterministic Model

- In 1996, Maurer released his *deterministic model*
- 4 statements: Authenticity, Trust, Recommendation, Certificate
- 2 inference rules:
  - ★ Derive authenticity
  - ★ Derive trust
- Initial View is the set of beliefs and observable statements
- Derived View is the initial view closed under inference rules
- If *Aut* is in my derived view, I can use the public key

# The Limits of Maurer's Model

- Authenticity of public keys
- Names = limited applicability
- Recommendation = all-or-none
- No time = no revocation or past
- No verification = bad deductions

# The Limits of Maurer's Model

- Authenticity of public keys → Binding b/t key and cert info
- Names = limited applicability → Properties, maybe name
- Recommendation = all-or-none → Trust transfer of properties
- No time = no revocation or past → Added time
- No verification = bad deductions → Added validity templates

# Definition 1: Statements

- **Authenticity of binding:**  $Aut(A, X, \mathcal{P}, \mathcal{I}) \stackrel{def}{=} A \xrightarrow{\mathcal{P}, \mathcal{I}} X$
- **Trust:**  $Trust(A, X, \mathcal{D}, \mathcal{I}) \stackrel{def}{=} A \xrightarrow{\mathcal{D}, \mathcal{I}} X$
- **Certificates:**  $Cert(X, B, \mathcal{P}, \mathcal{I}) \stackrel{def}{=} X \xrightarrow{\mathcal{P}, \mathcal{I}} B$
- **Trust Transfers:**  $Tran(X, Y, \mathcal{P}, \mathcal{I}) \stackrel{def}{=} X \xrightarrow{\mathcal{P}, \mathcal{I}} Y$
- We added second-order structures
  - ★ **Certificate Validity Templates:**  $Valid\langle A, Cert, t \rangle$
  - ★ **Transfer Validity Templates:**  $Valid\langle A, Tran, t \rangle$

## Definition 2: Inference Rules

- $View_A$  is Alice's initial view
- $\overline{View_A(t)}$  is Alice's derived view at time  $t$  where:

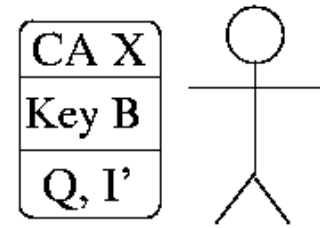
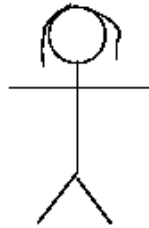
$\forall X, Y, t \in \{\mathcal{I}_0 \cap \mathcal{I}_1\}, Q \subseteq \mathcal{D} :$

$Aut(A, X, \mathcal{P}, \mathcal{I}_0), Trust(A, X, \mathcal{D}, \mathcal{I}_1), Valid\langle A, Tran(X, Y, Q, \mathcal{I}_2), t \rangle$   
 $\vdash Trust(A, Y, Q, \mathcal{I}_2)$

$Aut(A, X, \mathcal{P}, \mathcal{I}_0), Trust(A, X, \mathcal{D}, \mathcal{I}_1), Valid\langle A, Cert(X, B, Q, \mathcal{I}_2), t \rangle$   
 $\vdash Aut(A, B, Q, \mathcal{I}_2)$

- For  $A$  to believe  $B$  at time  $t$ ,  $Aut(A, B, Q, \mathcal{I}_2) \in \overline{View_A(t)}$

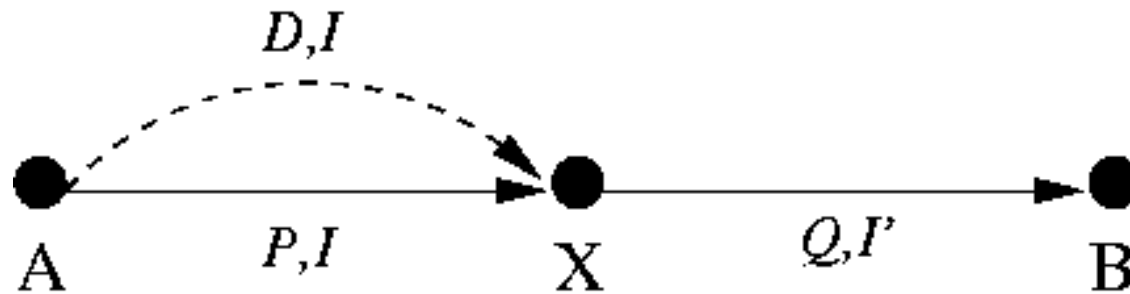
# An Example



- Alice and Bob both use CA  $X$
- $X$  certified Bob and assigned him properties  $Q$  for  $\mathcal{I}'$

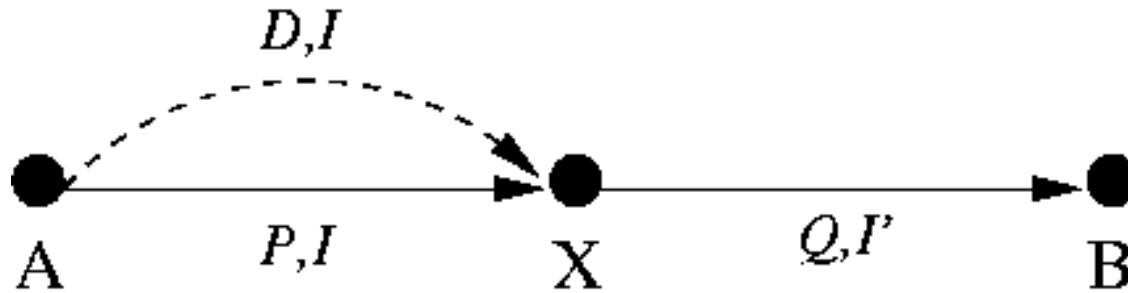


# Statement Graph



- $View_A = \{Aut(A, X, P, I), Trust(A, X, D, I), Cert(X, B, Q, I')\}$

# Statement Graph



- $View_A = \{Aut(A, X, P, I), Trust(A, X, D, I), Cert(X, B, Q, I')\}$
- Using the inference rules:  
$$Aut(A, X, P, I), Trust(A, X, D, I), Valid(A, Cert(X, B, Q, I'), t) \vdash Aut(A, B, Q, I')$$
- $\overline{View_A(t)} = View_A \cup Aut(A, B, Q, I')$
- Since  $Aut(A, B, Q, I') \in \overline{View_A(t)}$ , Alice uses Bob's cert

# Using the New Model

- Properties allow multiple cert families: X.509, ACs, PCs, SDSI/SPKI
- Time allows revocation and events in the past/future
- Properties allow for authorization scenarios
- Trust Transfers and domains enable delegation
- Time and Properties allow us to model hybrid PKIs: Greenpass and MyProxy

# Conclusions and Future Work

- New model can handle many types of real-world systems
- How well do the cert properties match the real world?
- Nonmonotonicity: decoupling cert lifespans from beliefs
- What kind of set operations on properties should we allow?